

Bratteli-Vershik Models of Cantor Minimal Systems

Maryam Hosseini

CIMPA School, Bejaia, Algeria

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there exists a Bratteli-Vershik system (X_B, T_B) that is conjugate to that.

To show this we need **Kakutani-Rokhlin partitions** that are based on the notion of **first return time to a clopen set**.

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$n(x)$ exists for every x as the system is minimal. In fact, the above sets have bounded gaps. This was because:

$$\forall \text{ clopen } U \subset X, \exists N; \bigcup_{n=0}^N T^{-n}(U) = X.$$

Example, return words for subshifts.

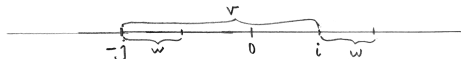
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$$\forall x \exists i; x_{[i, i+|w|-1]} = w \Rightarrow \exists j > 0; x_{[-j, i-1]} \in \mathcal{R}(w).$$

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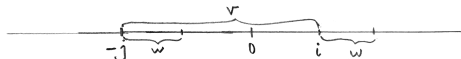


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So $S^i(x), S^{-j}(x) \in U$ and one of the the return times is:

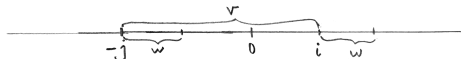
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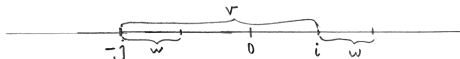
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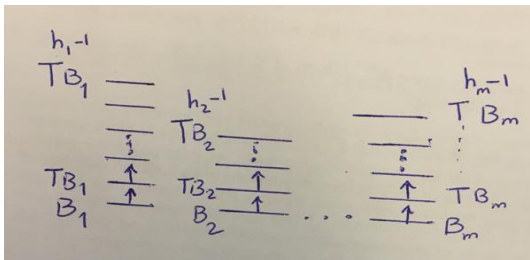
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So for $U = [w]$ and x :

$$n(x) = \min_{1 \leq i \leq m} |v_i|.$$

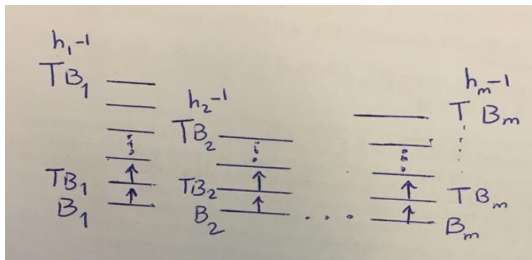
Kakutani-Rokhlin partitions.

Let (X, T) be a minimal Cantor system where T is a **homeomorphism**. A **Kakutani-Rokhlin** partition, say \mathcal{P} , of (X, T) is a collection of towers of the form:



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In fact,

$$\mathcal{P} = \{T^j(B_i) : 1 \leq i \leq m, 1 \leq j \leq h_i\}$$

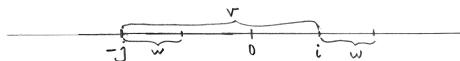
where $\{B_i\}_{i=1}^m$ are **disjoint clopen subsets** of X and $h_i \in \mathbb{N}$.

- $B = \cup_{i=1}^m B_i$ is called **the basis of partition**.
- for each $1 \leq i \leq m$, $\{B_i, TB_i, \dots, T^{h_i} B_i\}$ is called **a tower**. So the partition is a finite union of towers.
- For each $1 \leq i \leq m$ the **first return time to B** for the points of B_i is h_i
- The map T is defined explicitly every where **except the top levels** of the towers. We just know that the top level will be mapped to B but where exactly? we don't know.

Example, using return words.

Consider (X, S) as a minimal two-sided subshift, a word w and $U = [w] \subset X$. By minimality

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In fact, by minimality the set of return words to w is finite:

$$\mathcal{R}(w) = \{v_1, v_2, \dots, v_m\} : m < \infty.$$

So for each $v_i \in \mathcal{R}(w)$, $i = 1, \dots, m$ we have a tower above that.

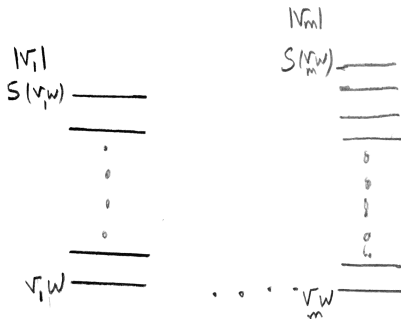
So let $B = \cup_{i=1}^m [V_i]$,

$$\mathcal{T}_i = \{S^j([v_i]) : v_i \in \mathcal{R}(w), 0 \leq j \leq |v_i|\}, i = 1, \dots, m$$

are the towers and

$$B = \{\mathcal{T}_i\}_{i=1}^m$$

is a Kakutani-Rokhlin partition for (X, S) .



Example.

Consider the two sided subshift $(\overline{\mathcal{O}(x)}, S)$ generated by the primitive substitution:

$$\sigma : 0 \mapsto 001, \quad 1 \mapsto 01$$

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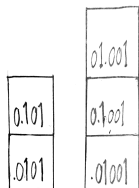
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Let $U = [w]$, with $w = 01$, as a clopen set around the fixed point of the substitution: Then

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Therefore, $\mathcal{B} = \{[0101], S[(0101)], [01001], S([01001]), S^2([01001])\}$ is a K-R partition of (X, S) .



Nested K-R partitions.

The K-R Partition \mathcal{B}' with basis B' is called **nested in** the K-R partition \mathcal{B} with basis B if $B' \subset B$ and \mathcal{B}' is a **refinement** of \mathcal{B} , i.e.

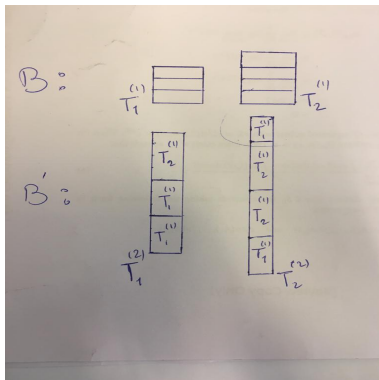
$$\forall A \in \mathcal{B}' \exists B \in \mathcal{B}; A \subset B.$$

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$$\forall A \in \mathcal{B}' \exists B \in \mathcal{B}; A \subset B.$$

Roughly speaking, we **cut and stack** the towers of \mathcal{B} to create \mathcal{B}' .



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- The sequence $\{\mathcal{B}^{(n)}\}_{n \geq 0}$ is nested, i.e. $\mathcal{B}^{(n+1)}$ is nested in $\mathcal{B}^{(n)}$ for every n .
- $\bigcup_n \mathcal{B}^{(n)}$ generates the topology of X .

K-R partitions of minimal Cantor systems.

Proposition

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$$\forall x \in C : n_C : X \rightarrow \mathbb{Z}; n_C(x) = \inf\{n > 0 : T^n x \in C\}.$$

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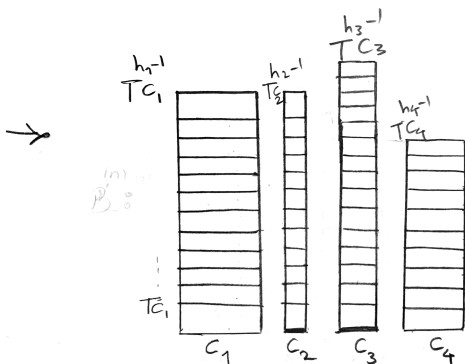
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Since (X, T) is minimal, n_C is well-defined and continuous and in fact,

$$C = \bigcup_{i=1}^m C_i; \quad \forall x \in C_i : n_C(x) = h_i \in \mathbb{N}.$$

Let's define a partition by the above subsets C_i , $i = 1, \dots, m$.

$$\mathcal{B}^{(1)} := \{T^j C_i : 0 \leq j \leq h_i - 1, 1 \leq i \leq m\}.$$



- every two cells in $\mathcal{B}^{(1)}$ are disjoint. This is just because of definition of n_C and T being a homeomorphism. (exercise)

- $\mathcal{B}^{(1)}$ covers X . Because

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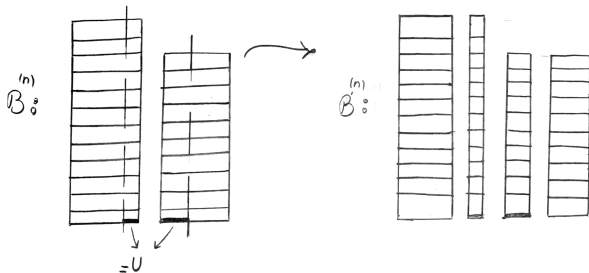
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Suppos that there exists some $U \in \mathcal{P}_1$ so that it has intersection with C_{ij} but C_{ij} is not contained in U . As the number of elements in $\mathcal{B}^{(1)}$ is finite, consider the least diameter δ of intersections of C_{ij} 's with the elements of \mathcal{P}_1 and cut the towers of $\mathcal{B}^{(1)}$ into finitely many thinner towers (with diameter δ) that each cell of these new towers are contained in an element of \mathcal{P}_1 .

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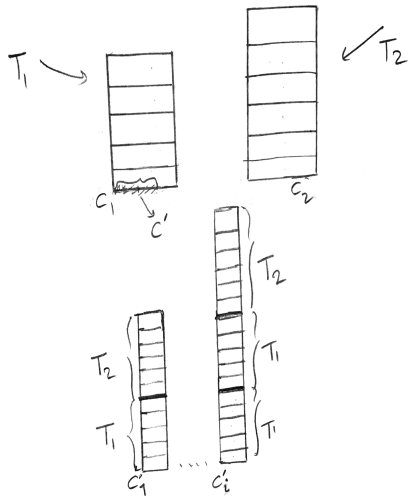
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By similar arguments as for $\mathcal{B}^{(1)}$, one can deduce that $\mathcal{B}^{(2)}$ is a K-R partition. Now we show that this is finer than $\mathcal{B}^{(1)}$ and could be constructed to be finer than \mathcal{P}_2 as well.

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- About $\mathcal{B}^{(1)}$: Since $C' \subset C_1$, the initial h_1 levels of all the towers in $\mathcal{B}^{(1)}$ are arranged in the same shape as the first tower of $\mathcal{B}^{(1)}$ with base C_1 . And since $T^{h_1}(C') \subset C_i$ for some other $i = 1, \dots, m$ the next h_i levels of each tower of $\mathcal{B}^{(2)}$ is arranged in the same way as the associated i 'th tower in $\mathcal{B}^{(1)}$ and in fact, each of these levels is contained in a unique level of the associated i 'th tower in $\mathcal{B}^{(1)}$.



By continuing the above procedure inductively by choosing nested sequence of clopen subsets of C_1 around x_0 and constructing towers based on them and so a sequence of refining K-R partitions that each of them refines a new \mathcal{P}_n (increasingly by n), we will have the intersection of all the bases of the towers to be $\{x_0\}$.

Example.

Let's get back to the primitive substitution:

$$\sigma : 0 \mapsto 001, \quad 1 \mapsto 01$$

and $C = [w]$, $w = 0$, as a clopen set around the fixed point of the substitution:

$$x = 01001010010010100101 \dots$$

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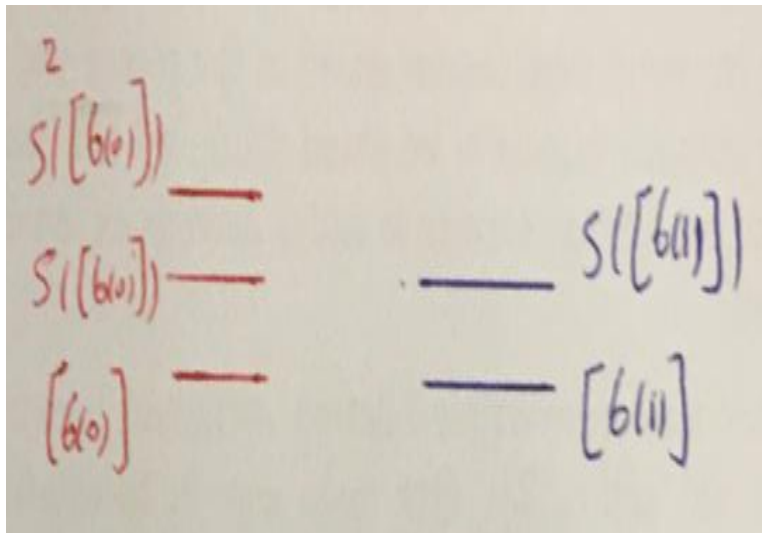
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Then

$$\mathcal{R}(w) = \{0, 01\}, \quad [0] = [00] \cup [01] = [001] \cup [01] = [\sigma(0)] \cup [\sigma(1)].$$

So we should have two towers with disjoint bases $[001]$ and $[01]$. So

$$\begin{aligned} \mathcal{B}^{(1)} &= \{[\sigma(0)], S([\sigma(0)]), S^2([\sigma(0)]), [\sigma(1)], S([\sigma(1)])\} \\ &= \{T_1^{(1)}, T_2^{(1)}\} \end{aligned}$$



Now let $C'_1 = [w']$, $w' = 001$. Then

$$\mathcal{R}(w) = \{001, 00101\}, \quad [001] = [00100101] \cup [00101] = [\sigma^2(0)] \cup [\sigma^2(1)]$$

So we should have two towers with disjoint bases $[001001]$ and $[00101]$:

$$\begin{aligned} \mathcal{B}^{(2)} &= \{[\sigma^2(0)], \dots, S^7([\sigma^2(0)]), [\sigma^2(1)], \dots, S^4([\sigma^2(1)])\} \\ &= \{T_1^{(1)}, T_1^{(1)}, T_2^{(1)}, T_1^{(1)}, T_2^{(1)}\} \\ &= \{T_1^{(2)}, T_2^{(2)}\} \end{aligned}$$

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If you continue this procedure inductively, you will get

$$\begin{aligned} \mathcal{B}^{(n)} &= \{S^j \sigma^n([a]) : a = 0, 1, 0 \leq j < |\sigma^n(a)|\} \\ &= \{T_1^{(n-1)}, T_1^{(n-1)}, T_2^{(n-1)}, T_1^{(n-1)}, T_2^{(n-1)}\} \\ &= \{T_1^{(n)}, T_2^{(n)}\} \end{aligned}$$

$S^2([6_{(0)}])$ —
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$|6_{(0)}| - 1$
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General case.

Proposition

For every aperiodic two sided minimal subshift (X, S) associated to a primitive proper substitution $\sigma : A \rightarrow A^*$ that A is a finite alphabet,

$$\mathcal{B}^{(n)} = \{S^j \sigma^n([a]) : a \in A, 0 \leq j < |\sigma^n(a)|\}, \quad n \geq 0$$

is a refining sequence of K - R partitions.

Remark. Proper means that there are letters $r, \ell \in A$ that for every $a \in A$, $\sigma(a)$ starts with r and ends up with ℓ .

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To prove this proposition knowing the notion of recognizability will ease it.

